

Paolo Lucente

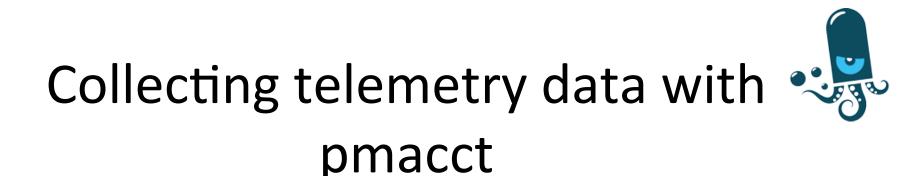
pmacct

AfPIF 2017, Abidjan – Aug 2017



Presentation history

- **1**.1:
 - MENOG 13 meeting, Kuwait City, Sep 2013
- **1**.2:
 - SEE 3 meeting, Sofia, Apr 2014
- **1**.3:
 - AfPIF 2017 meeting, Abidjan, Aug 2017



Agenda

Introduction

- pmacct architecture & benefits
- example, data aggregation: traffic matrices
- example, logging micro-flows or events
- tee: briefly on horizontal scalability

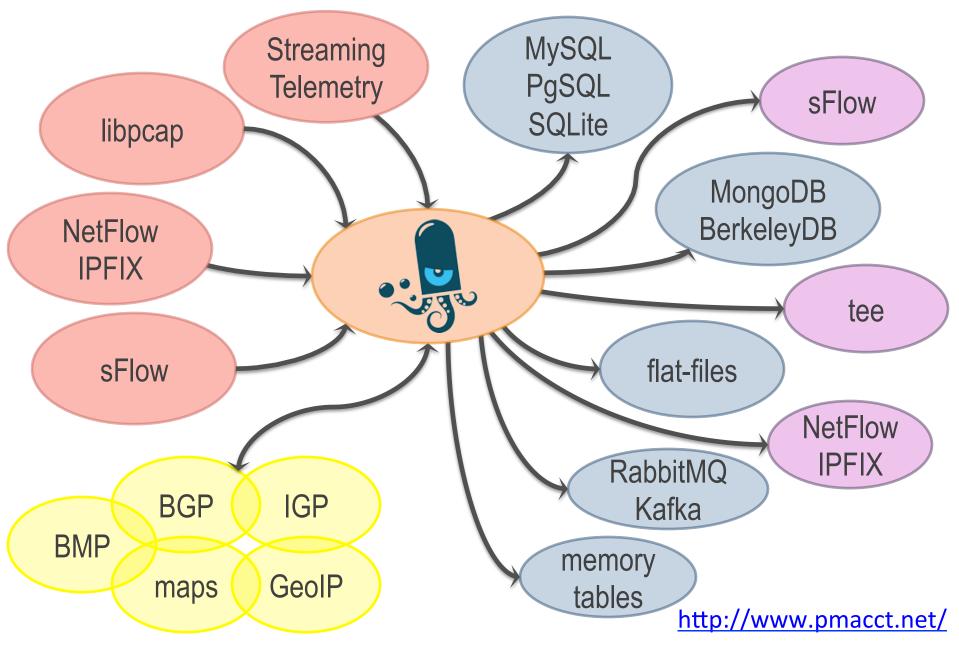
whoami

Paolo Lucente GitHub: <u>paololucente</u> LinkedIn: <u>plucente</u>

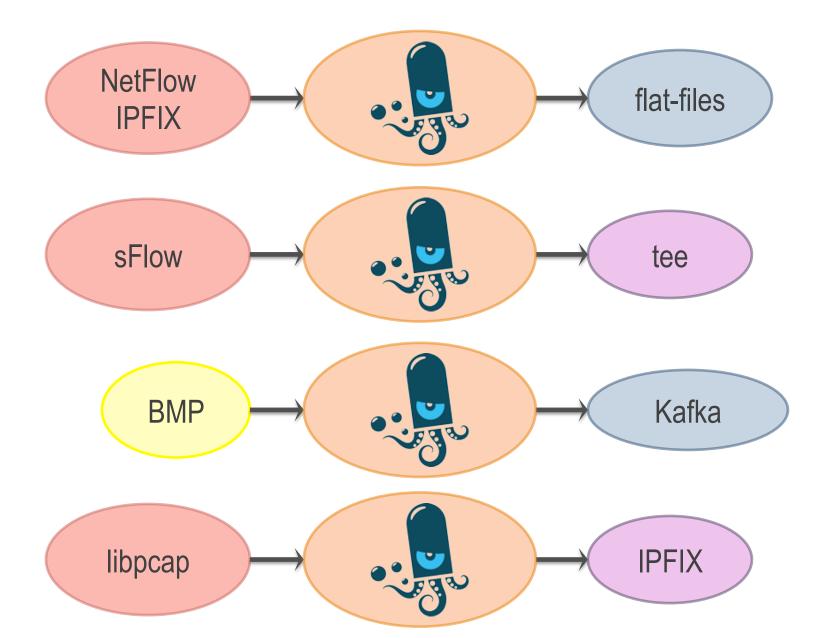


Digging data out of networks worldwide for fun and profit for more than 10 years

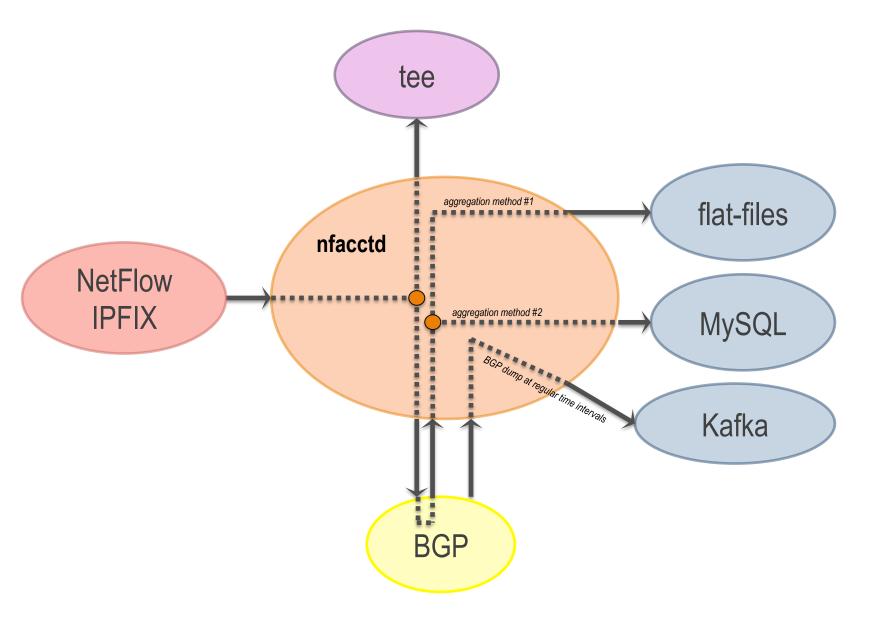
pmacct is open-source, free, GPL'ed software



pmacct: a few simple use-cases

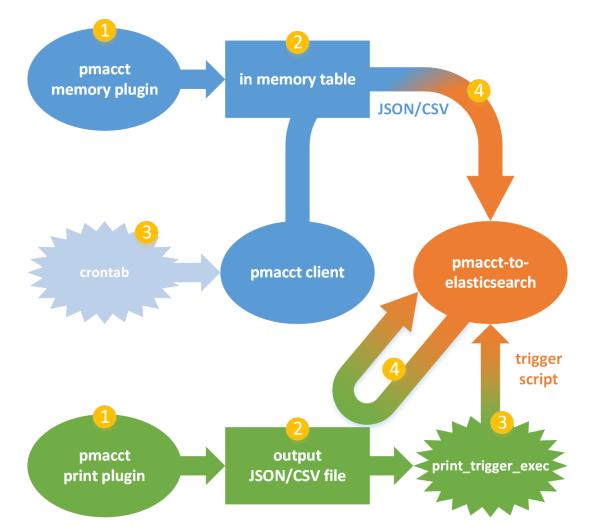


pmacct: a slightly more complex use-case





pmacct-to-elasticsearch 0.3.0



Credits to: Pier Carlo Chiodi, https://github.com/pierky/pmacct-to-elasticsearch

Use cases by industry

ISPs, Hotspots, Data-center

Monitor customer quotas or fair-usage policy Peering

IXPs

Infer member relations Provide members traffic stats Capacity planning Triggering alarms Historical traffic trends Feeding into 3rd party tools

Mobile operators

Verify roaming charges Inspect subscribers behaviour

IP Carriers, CDNs

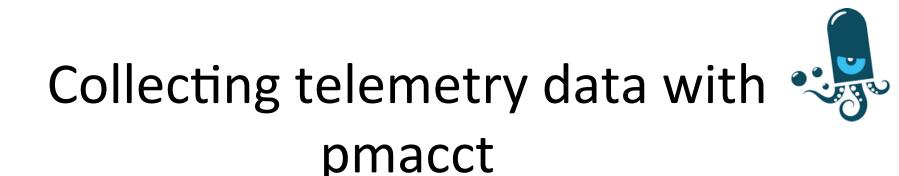
Detect revenue leaks Customer retention Peering

SDN

Query of traffic stats on custom spatial and temporal bounds

Key pmacct non-technical facts

- 10+ years old project
- Can't spell the name after the second drink
- Free, open-source, independent
- Under active development
- Innovation being introduced
- Well deployed around, also in large SPs/IXPs
- Close to the SP/IXP community needs



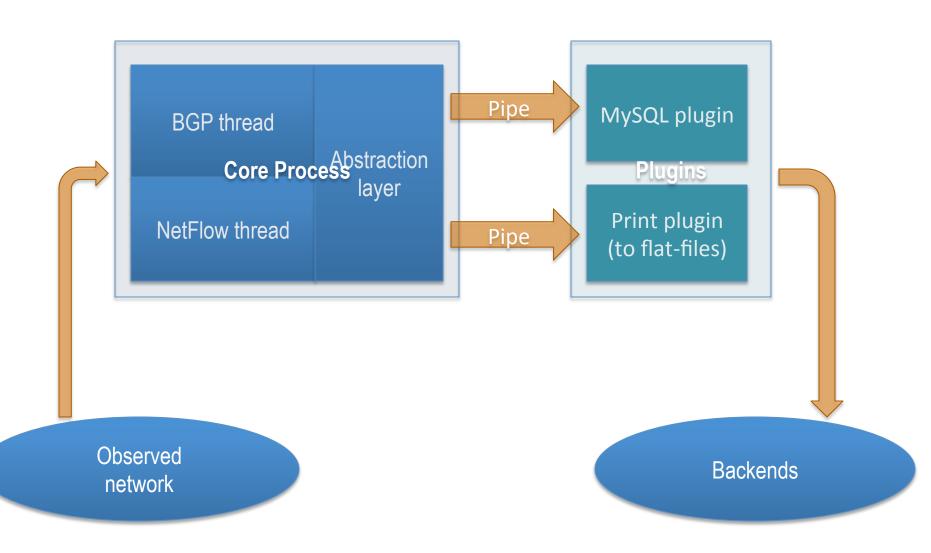
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Some technical facts

- Pluggable architecture:
 - Can easily add support for new data sources and backends
- Correlation of data sources:
 - Natively supported data sources (ie. flow telemetry, BGP, BMP, IGP, Streaming Telemetry)
 - Enrich with external data sources via tags and labels
- Enable analytics against each data source:
 - Stream real-time
 - Dump at regular time intervals (possible state compression)

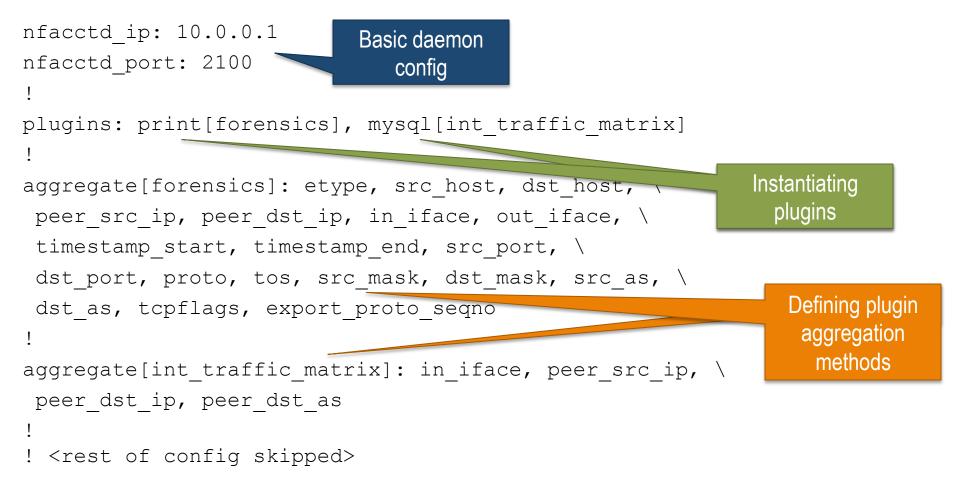
Some technical facts (cont.d)



Some technical facts (cont.d)

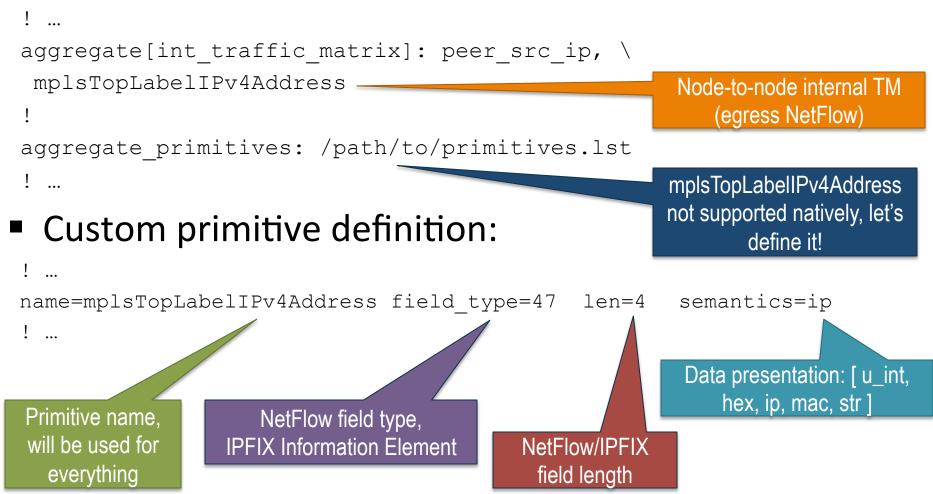
- Build multiple views out of the very same collected network traffic, ie.:
 - Unaggregated to flat-files for security and forensics; or to message brokers (RabbitMQ, Kafka) for Big Data
 - Aggregated as [<ingress router>, <ingress interface>, <BGP next-hop>, <peer destination ASN>] and sent to a SQL DB to build an internal traffic matrix for capacity planning purposes
- Pervasive data-reduction techniques, ie.:
 - Data aggregation
 - Filtering
 - Sampling

Touching ground: a config snippet for both aggregated and unaggregated views



Touching ground: data aggregation & custom-defined primitives

Config file snippet:



BGP integration

- pmacct introduced a Quagga-based BGP daemon:
 - Implemented as a parallel thread within the collector
 - Doesn't send UPDATEs whatsoever
 - Behaves as a passive BGP neighbor
 - Maintains per-peer BGP RIBs
 - Supports 32-bit ASNs; IPv4, IPv6 and VPN families
 - Supports ADD-PATH: draft-ietf-idr-add-paths
- Why BGP at the collector?
 - Telemetry reports on forwarding-plane, and a bit more
 - Extended visibility into control-plane information

BMP integration

- pmacct introduced a BMP daemon written from scratch:
 - BMP is: BGP Monitoring Protocol
 - Implemented as a parallel thread within the collector
 - All goodies already described for the BGP daemon:
 - Contributing to IETF draft-ietf-grow-bmp-local-rib
 - Visibility in Adj-RIB-In
 - Visibility in Adj-RIB-Out:
 - Contributing to IETF draft-ietf-grow-bmp-adj-rib-out

IGP (IS-IS) integration

- A Quagga-based IS-IS daemon was introduced:
 - Implemented as a parallel thread within the collector
 - IS-IS neighborship over a GRE tunnel
 - Currently limited to single neighborhip, single level, single topology
 - Useful to look up non BGP-routed networks
- It will get eventually replaced (BGP-LS)

Storing data persistently

- Data need to be aggregated both in spatial and temporal dimensions before being written down:
 - Optimal usage of system resources
 - Avoids expensive consolidation of micro-flows
- Build project-driven data set(s):
 - No shame in multiple partly overlapping data-sets
 - Optimize computing

Storing data persistently (cont.d)

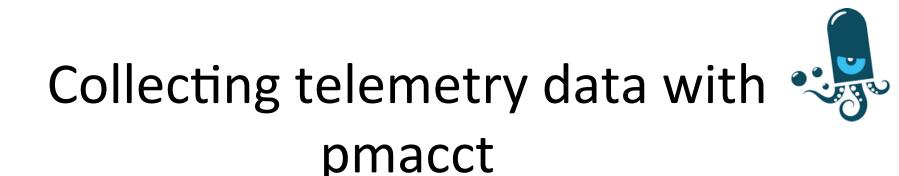
- "noSQL" databases (Big Data ^(C)):
 - Able to handle large time-series data-sets
 - Meaningful subset of SQL query language
 - Innovative storage and indexing engines
 - Scalable: clustering, spatial and temporal partitioning
 - UI-ready: ie. ELK and TICK stacks
- Open-source RDBMS:
 - Able to handle large data-sets
 - Flexible and standardized SQL query language
 - Solid storage and indexing engines
 - Scalable: clustering, spatial and temporal partitioning

Storing data persisently: MongoDB

- Once it was pmacct opening to noSQL databases:
 - Mess up with the C API over time
 - 2017 reality check: lack of interest, discontinuing
- noSQL landscape difficult to move through, ie.
 fragmented and lacks of standardization
- MongoDB seemed interesting for:
 - Native grouping operation (more performing and less complex than map/reduce)
 - Horizontal scaling concept (sharding)

Brokering data around: RabbitMQ, Kafka message exchanges

- noSQL landscape difficult to move through, ie. fragmented and lacks of standardization, am i repeating myself? ⁽ⁱ⁾
- Data can be picked up at the message exchange in the preferred programming/scripting language
- Data can be then easily inserted in the preferred backend, ie. not natively supported by pmacct



Agenda

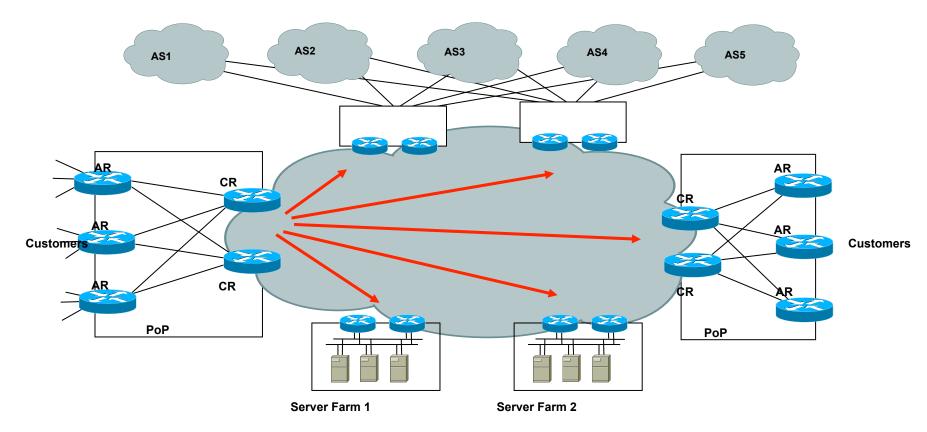
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- tee: briefly on horizontal scalability

Why speaking of traffic matrices?

- Are traffic matrices useful to a network operator in the first place? Yes ...
 - Capacity planning (build capacity where needed)
 - Traffic Engineering (steer traffic where capacity is available)
 - Better understand traffic patterns (what to expect, without a crystal ball)
 - Support peering decisions (traffic insight, traffic engineering at the border, support what if scenarios)

Review of traffic matrices: internal

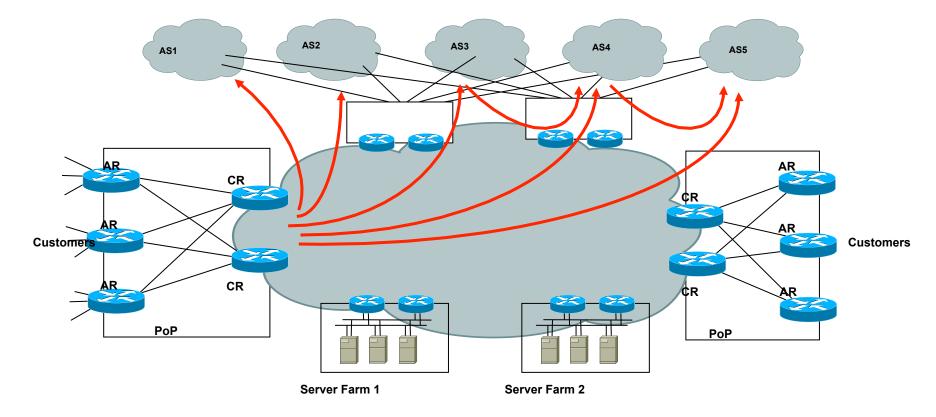
POP to POP, AR to AR, CR to CR



© 2010 Cisco Systems, Inc./Cariden Technologies, Inc.

Review of traffic matrices: external

 Router (AR or CR) to external AS or external AS to external AS (for IP transit providers)

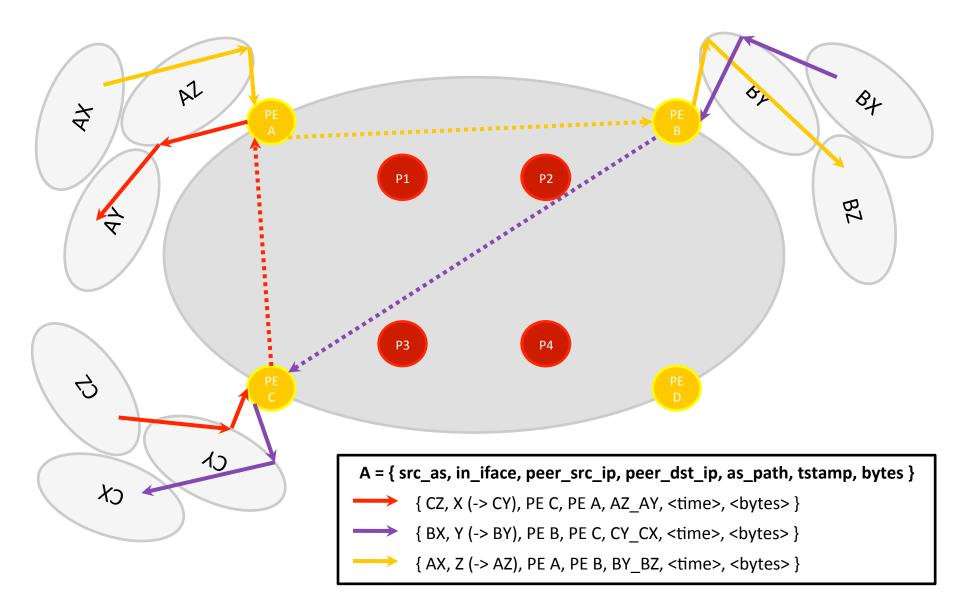


© 2010 Cisco Systems, Inc./Cariden Technologies, Inc.

Let's focus on an external traffic matrix to support peering decisions

- Analysis of existing peers and interconnects:
 - Support policy and routing changes
 - Fine-grained accounting of traffic volumes and ratios
 - Determine backbone costs associated to peering
 - Determine revenue leaks
- Planning of new peers and interconnects:
 - Who to peer next
 - Where to place next interconnect
 - Modeling and forecasting

Our traffic matrix visualized

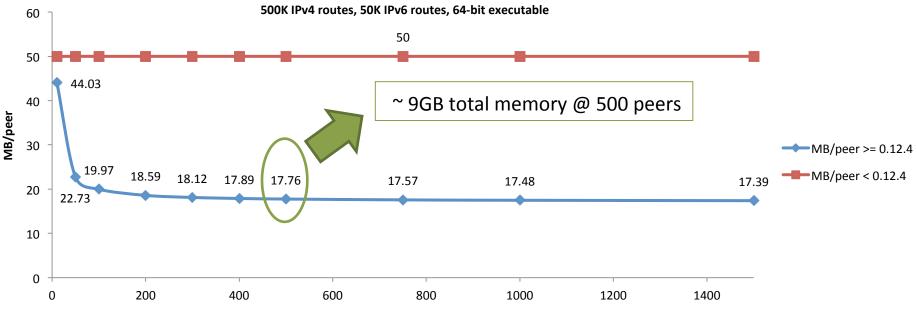


Getting BGP to the collector

- Needed for technical reasons:
 - Flow exporters use NetFlow v5, ie. no BGP next-hop
 - Flow exporters are unaware of BGP
 - Libpcap is used to collect traffic data
- Needed for topology or traffic related reasons:
 - Transiting traffic to 3rd parties
 - Dominated by outbound traffic

Getting BGP to the collector (cont.d)

- Let the collector BGP peer with all PE devices: facing peers, transit and customers.
- Determine memory footprint (below in MB/peer, using BGP best-path sessions)



Number of BGP peers

Getting BGP to the collector (cont.d)

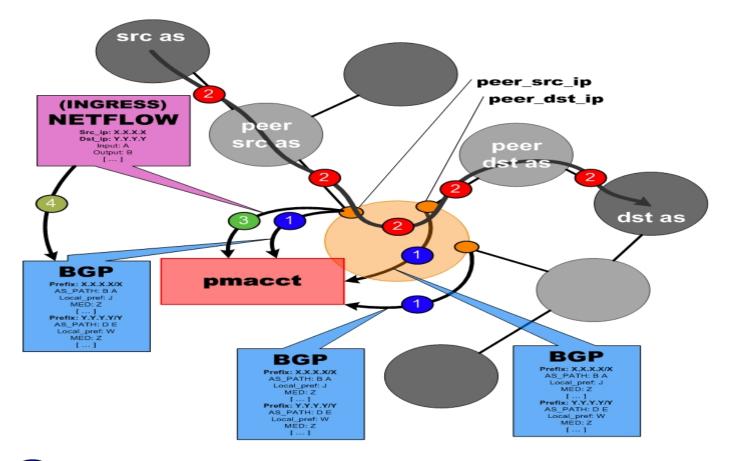
- Set the collector as iBGP peer at the PE devices:
 - Configure it as a RR client to get full table
 - Collector acts as iBGP peer across (sub-)ASes
- BGP next-hop has to represent the remote edge of the network model:
 - Typical scenario for MPLS networks
 - Can be followed up to cover specific scenarios like:
 - BGP confederations
 - default gateway defined due to partial or default-only routing tables

Getting telemetry to the collector

- Export ingress-only measurements at all PE devices: facing peers, transit and customers.
 - Traffic is routed to destination, so plenty of information on where it's going to
 - True, some eBGP multi-path scenarios may get challenging
 - It's crucial instead to get as much as possible about where traffic is coming from, ie.:
 - input interface at ingress router
 - source MAC address

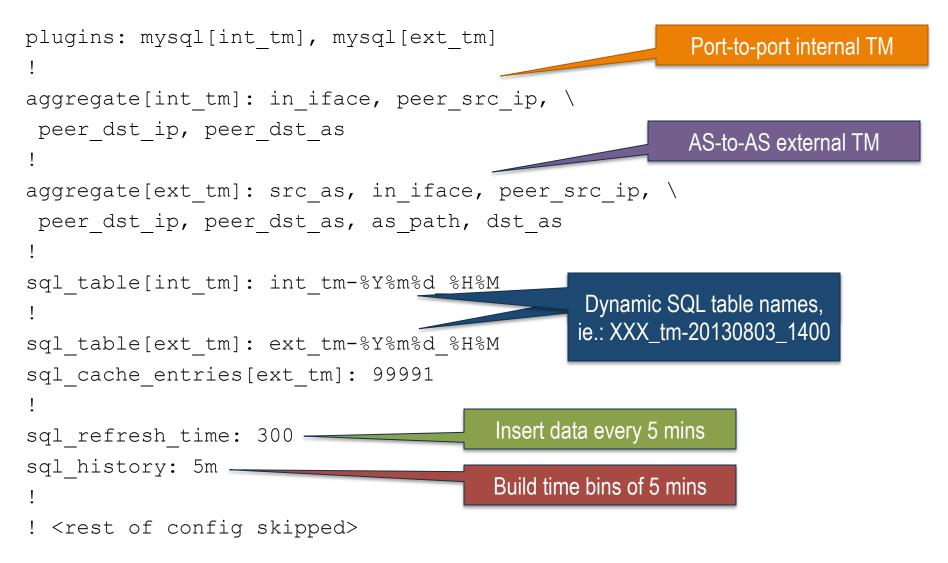
- Perform data reduction at the PE (ie. sampling)

Telemetry data/BGP correlation

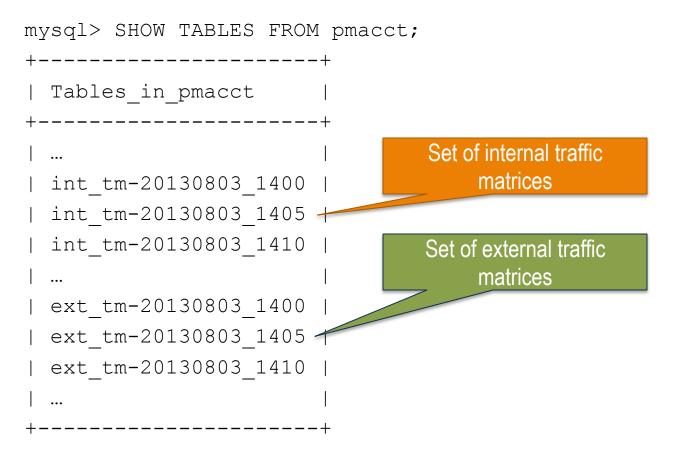


- Edge routers send full BGP tables to pmacct
- 2 Traffic flows
- NetFlow records are sent to pmacct
- Pmacct looks up BGP information: NF src addr == BGP src addr

Touching ground: a config snippet for traffic matrices



Touching ground: how data would look like: internal traffic matrix example (1/2)

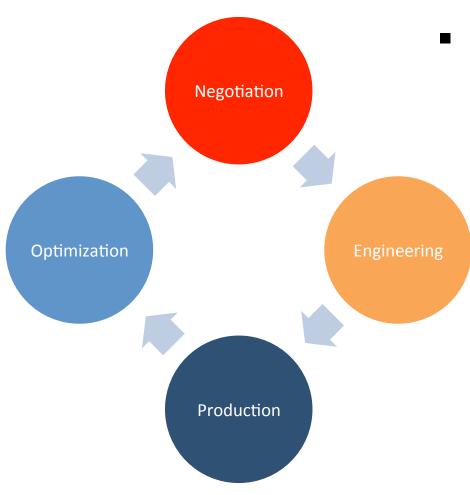


NOTE: sub-aggregation is expensive: we could also have had our traffic matrices over multiple temporal aggregations in parallel, ie. 5 mins (as above) but also hourly and daily.

Touching ground: how data would look like: internal traffic matrix example (2/2)

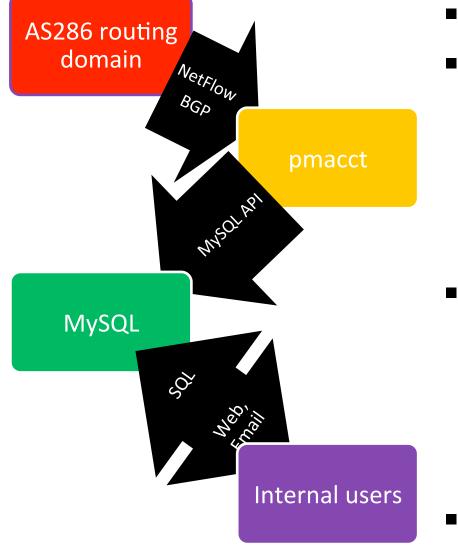
mysql> SELECT * FROM int tm-20130803 1400 LIMIT 10; iface in | peer ip src | peer ip dst | peer dst as | stamp inserted | bytes 212 | 10.0.0.107 | 10.0.0.3 65000 03-08-2013 14:00 859 212 | 10.0.0.107 | 10.0.0.253 03-08-2013 14:00 5358 65001 10.0.0.107 212 | 10.0.0.234 65002 03-08-2013 14:00 6181 03-08-2013 14:00 10.0.0.107 | 10.0.0.251 212 | 65003 27002 10.0.0.107 | 10.0.0.233 03-08-2013 14:00 205 L 65004 1200 | 10.0.0.240 258 | 10.0.0.107 65005 03-08-2013 14:00 560 10.0.0.107 10.0.252 65006 03-08-2013 14:00 62682 212 | 10.0.0.107 10.0.234 03-08-2013 14:00 212 65007 3843 10.0.0.107 10.0.0.17 03-08-2013 14:00 212 | 65008 21074 10.0.0.107 10.0.254 03-08-2013 14:00 2023 205 65009 10 rows n set / U Here is our matrix Getting ingress NetFlow Amount of traffic sent in the Time reference from 10.0.0.7 time window 14:00 – 14:05

Case-study: peering at AS286 [year 2009]



- Peering as a cycle
 - NetFlow + BGP traffic matrix steers peering optimization:
 - Identify new and "old" peers
 - Traffic analysis: backbone costs, 95th percentiles, ratios
 - Analysis of interconnection density and traffic dispersion
 - Forecasting and trending
 - Ad-hoc queries from Design & Engineering and indeed ... the IPT Product Manager

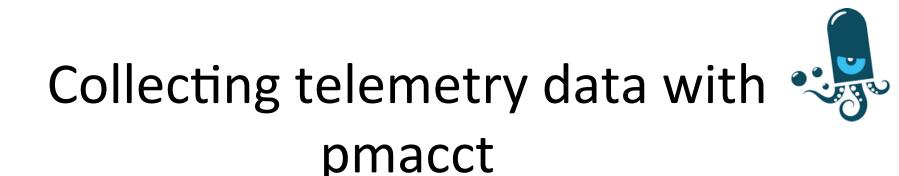
Case-study: peering at AS286 [year 2009]



- 250+ Gbps routing-domain
- 100+ high-end routers around the globe:
 - Export sampled NetFlow
 - Advertise full routing table
 - Mix of Juniper and Cisco
- Collector environment:
 - Runs on 2 Solaris/SPARC zones
 - pmacct: dual-core, 4GB RAM
 - MySQL: quad-core, 24GB RAM, 500 GB disk
 - Data retention: 6 months

Case-study: peering at AS286 [year 2009]

- AS286 backbone routers are first configured from templates:
 - NetFlow + BGP collector IP address defined over there
 - Enabler for auto-discovery of new devices
- Edge interfaces are provisioned following service delivery manuals:
 - Relevant manuals and TSDs include NetFlow activation
 - Periodic checks NetFlow is active where it should
- Maps, ie. source peer-AS to ingress interfaces, are re-built periodically



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pmacct, logging & flat-files: brief history (1/2)

- Originally pmacct was about memory tables and RDBMS (no flat-files)
- It was also about data aggregation (no logging micro-flows)
- Other tools were doing this greatly already better invest in new ideas
- In 2011 pmacct opened to flat-files (see next slide)

pmacct, logging & flat-files: brief history (2/2)

- In recent years the landscape changed and NetFlow and IPFIX protocols were being generalized besides flows
- Noticeably, they entered in the Event Logging space, ie.:
 - Cisco NEL, ie. (CG)NAT events
 - Cisco NSEL, ie. security events
- This was great time to review the strategy around flat-files and logging (of events and, as a consequence, of micro-flows)

Logging use-cases

- Micro-flows:
 - R&D, lab activities
 - Security, DDoS detection, forensics
 - Related to data aggregation:
 - Analysis, temporary: elaborate aggregation methods
 - Back-up, permanent: keep raw data in case useful
- Events:
 - NAT events: compliancy with local authorities
 - FW events: security reports

Logging features

- Split data in files and directories basing on time of the day and (some) primitive values, ie.:
 - 5 mins, hourly, daily files (fully customizable)
 - IP address of telemetry exporter
- Can append to an existing file: ie. hourly files but refreshed every 5 mins
- Files can be then archived right away via triggers, print_trigger_exec)
- Pointer to latest file in time-series

Logging data formats

- Text formats supported: tab-spaced, CSV and JSON
 - CSV: quickest serialization, not self-descriptive
 - JSON: verbose, (hence) slower to serialize, selfdescriptive
- Binary format supported: Apache Avro
 - Space efficient, schema-based, indirect access
 - GPB being avoided
 - Maybe revisit Cap'n Proto in future

Touching ground: a config snippet for logging micro-flows

<pre>plugins: print[forensics]</pre>	Micro-flow (de)aggregation				
<pre>aggregate[forensics]: src_host, dst_host, \ peer_src_ip, peer_dst_ip, in_iface, out_iface, \ timestamp_start, timestamp_end, src_port, \</pre>					
<pre>dst_port, proto, tos, src_mask, dst_mask, src_as, dst_as, tcpflags</pre>	Dynamic file names, ie.: forensics-20130803_1400				
print_output_file[forensics]: /path/to/forensics-%Y%m%d_%H%M.txt					
<pre>print_output[forensics]: csv print_refresh_time[forensics]: 300 print_bistory[forensics]: 5m</pre>	[formatted, csv, json]				
<pre>print_history[forensics]: 5m print_output_file_append[forensics]: true !</pre>	Insert data every 5 mins, append to file if exists				
<pre>print_latest_file[forensics]: /path/to/forensics-l</pre>	latest				
! <rest config="" of="" skipped=""></rest>	Pointer to latest file to become optional and				

explicitly configured

Touching ground: how data would look like: logging micro-flows (1/2)

shell> ls -la

•••

-rw	1	pmacct	pmacct
-rw	1	pmacct	pmacct
-rw	1	pmacct	pmacct
-rw	1	pmacct	pmacct
lrwxrwxrwx forensics-2		-	-
101010100 2	с <u>т</u> .	5555 <u>5</u> 11	00.0110

<size> Aug 02 13:50 forensics-20130802-1345.txt
<size> Aug 02 13:55 forensics-20130802-1350.txt
<size> Aug 02 14:00 forensics-20130802-1355.txt
<size> Aug 02 14:05 forensics-20130802-1400.txt
10 Aug 02 14:05 forensics-latest -> /path/to/

Configurable ownership

Pointer to latest finalized file

Touching ground: how data would look like: logging micro-flows (2/2)

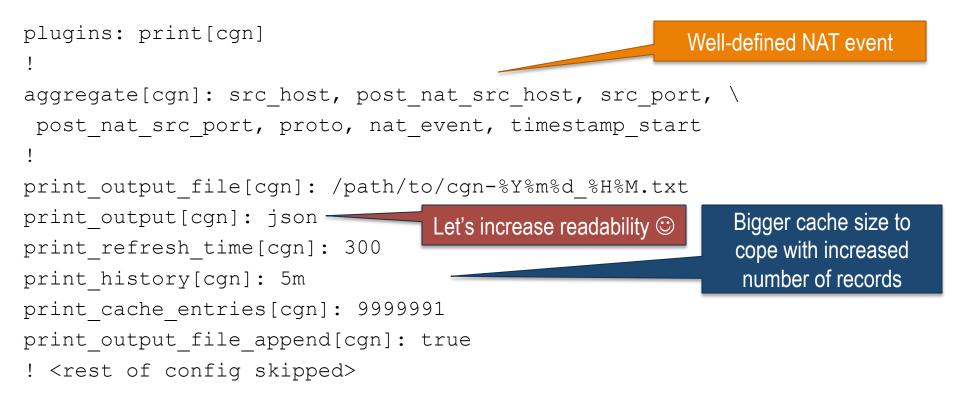
shell> cat forensics-latest

SRC_AS, DST_AS, PEER_SRC_IP, PEER_DST_IP, IN_IFACE, OUT_IFACE, SRC_IP, DST_IP, SRC_MASK, DST_MASK, SRC_PORT, DST_PORT, TCP_FLAGS, PROTOCOL, TOS, TIMESTAMP_ST ART, TIMESTAMP_END, PACKETS, BYTES 65001, 65002, 10.0.0.1, 10.0.0.100, 101, 8, 192.168.158.133, 192.168.126.141, 2 4, 24, 61912, 22, 24, tcp, 16, 2013-08-04 17:40:12.167216, 2013-08-04

17:41:36.140279,21,1407

[..]

Touching ground: a config snippet for logging NAT events



NOTE 1: see config snippet for micro-flows (a few slides back) for additional comments on configuration directives listed above.

NOTE 2: a bigger cache is beneficial to limit scattering of writes to the backend. If the configured cache is unable to contain all records, a purge of data to the backend is triggered and cache content is flushed so to make room to new data.

Touching ground: how data would look like: logging NAT events (1/2)

shell> ls -la

-rw	1	pmacct	pmacct
-rw	1	pmacct	pmacct
-rw	1	pmacct	pmacct
-rw	1	pmacct	pmacct
lrwxrwxrwx cgn-20130802		—	pmacct

<size></size>	Aug	02	13:50	cgn-20130802-1345.txt
<size></size>	Aug	02	13:55	cgn-20130802-1350.txt
<size></size>	Aug	02	14:00	cgn-20130802-1355.txt
<size></size>	Aug	02	14:05	cgn-20130802-1400.txt
10 Aug	02 1	4:0)5 cgn-	-latest -> /path/to/

•••

•••

Touching ground: how data would look like: logging NAT events (2/2)

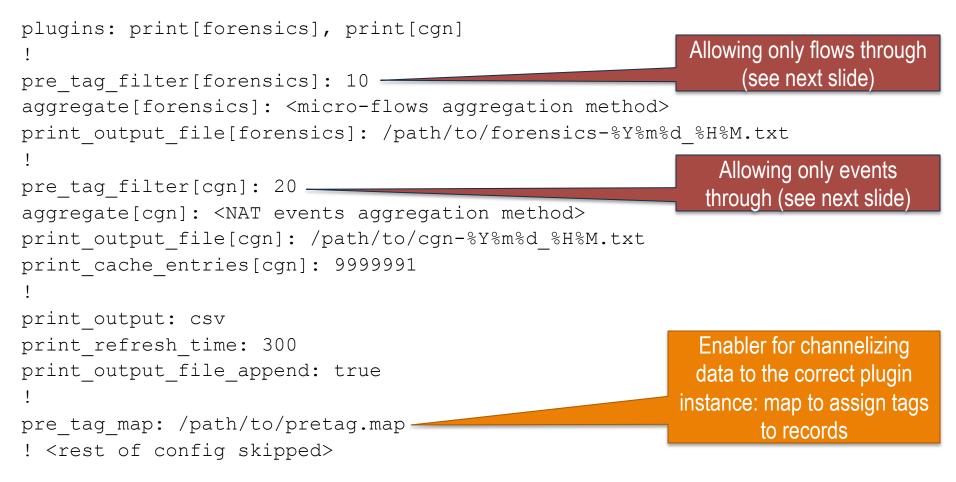
shell> cat cgn-latest

{"timestamp_start": "2013-02-21 16:56:33.518000000", "ip_proto": "tcp", "post_nat_ip_src": "1.2.179.16", "ip_src": "192.168.37.51", "port_src": 61020, "nat_event": 1, "post_nat_port_src": 31392} [..]

A single (set of) collector(s) for both micro-flows and events logging?

- Yes, possible:
 - All NetFlow, regardless, pointed to the same place
 - Makes sense on small-medium deployments
 - On larger ones potentially pressuring the (same set of) collector(s) with, say, an ongoing DDoS and a CGN blade rebooting is not a good idea. Go for splitting.
- pmacct able to tag and channelize data (ie. send data selectively to plugins) basing on a number of clauses

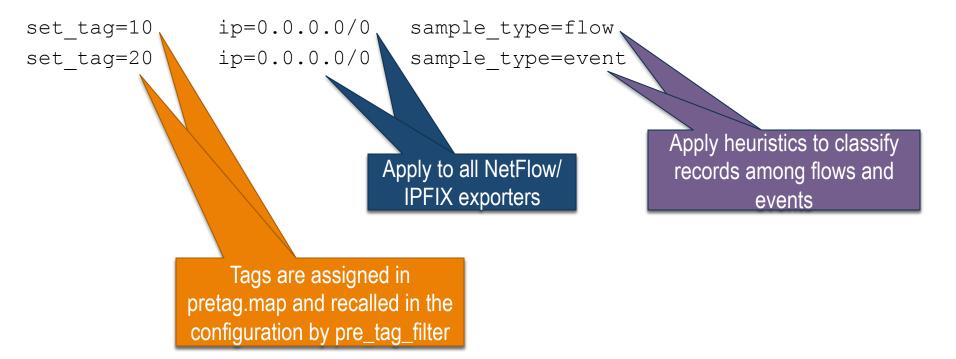
Touching ground: a config snippet for both micro-flows and event logging (1/2)

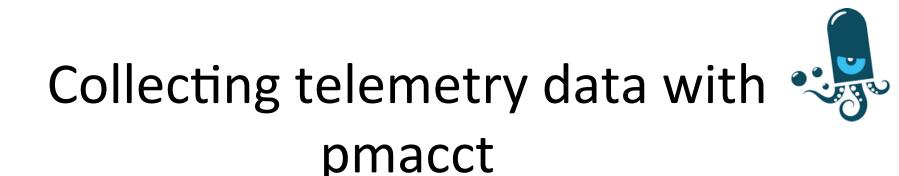


NOTE: This configuration merely merges the micro-flows and event logging configurations seen before. Check them out (a few slides back) for additional comments on configuration directives listed above.

Touching ground: a config snippet for both micro-flows and event logging (2/2)

shell> cat pretag.map





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Briefly on scalability

- A single collector might not fit it all:
 - Memory: can't store all BGP full routing tables
 - CPU: can't cope with the pace of telemetry export
- Divide-et-impera approach is valid:
 - Assign PEs (both telemetry and BGP) to collectors
 - If natively supported DB:
 - Assign collectors to DB nodes
 - Cluster the DB
 - If not-natively supported DB:
 - Assign collectors to message brokers
 - Cluster the messaging infrastructure

Briefly on scalability (cont.d)

- Intuitively, the matrix can become big:
 - Can be reduced by excluding entities negligible to the specific scenario:
 - Keep smaller routers out of the equation
 - Filter out specific (class of) customers
 - Focus on downstream if CDN, upstream if ISP
 - Sample or put thresholds on traffic relevance

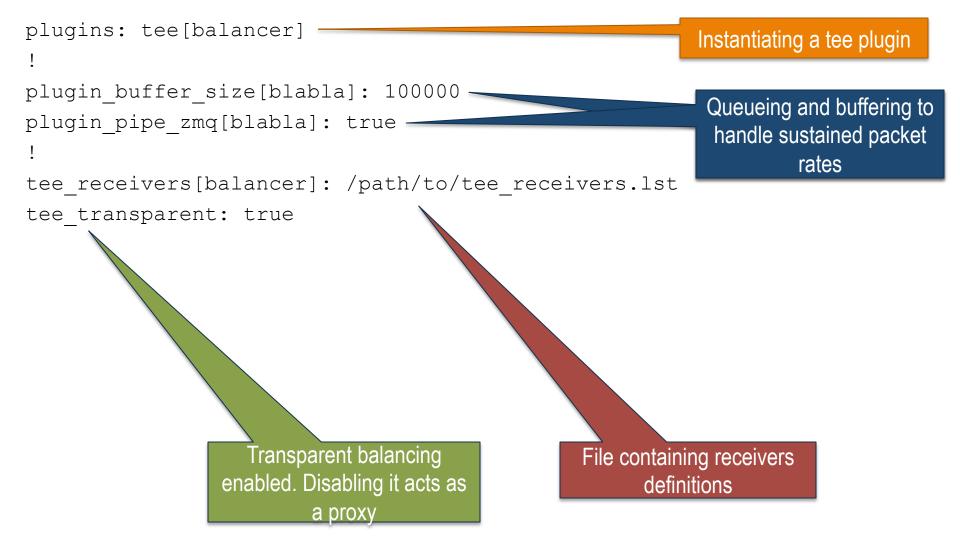
Need for horizontal scalability in a telemetry collector

- High-frequency sampling (ie. security)
- Cope with increasing data rates:
 - 10G to 100G but, depending on the application, sampling rates might stay the same
 - Events logging: ie. NetFlow challenges Syslog in the space of logging Carrier Grade NAT (CGN) and firewall events
- Scale without super-computing powers

pmacct & horizontal scaling

- Supports a 'tee' plugin
 - Receivers can be added/changed/removed on the fly
 - Load-balenced tee'ing (hashed or round-robin)
 - Selective tee'ing
- Multiple pmacct collectors can run in parallel
 - Coupling telemetry and routing data from same PE

Touching ground: a config snippet for transparent hashing balancer (1/2)



Touching ground: a config snippet for transparent hashing balancer (2/2)

shell> cat tee receivers.lst

id=1 \setminus

ip=192.168.5.1:2100,192.168.5.2:2100,192.168.5.3:2100 \
balance-alg=hash-agent

Touching ground: a config snippet for transparent selective balancer (1/2)

```
plugins: tee[balancer]
!
plugin_buffer_size[blabla]: 100000
plugin_pipe_zmq[blabla]: true
!
tee_receivers[balancer]: /path/to/tee_receivers.lst
tee_transparent: true
!
pre tag map: /path/to/pretag.map
```

Enabler for selective balancing: map to assign tags to NetFlow/IPFIX exporters

NOTE: see config snippet for transparent hashing balancer (a few slides back) for additional comments on configuration directives listed above.

Touching ground: a config snippet for transparent selective balancer (2/2)

shell> cat tee receivers.lst

- id=2 ip=192.168.4.1:2100 tag=100-
- id=3 ip=192.168.4.2:2100 tag=200

shell> cat pretag.map

Tags are assigned in pretag.map and recalled in tee_receivers.lst

set_tag=100 ip=10.0.0/25

set tag=200 ip=10.0.0.128/25

Further information about pmacct

- <u>https://github.com/pmacct/pmacct</u>
 - Official GitHub repository, where star and watch us $\ensuremath{\mathfrak{S}}$
- http://www.pmacct.net/lucente_pmacct_uknof14.pdf
 - More about coupling telemetry and BGP
- http://ripe61.ripe.net/presentations/156-ripe61-bcpplanning-and-te.pdf
 - More about traffic matrices, capacity planning & TE
- <u>https://github.com/pmacct/pmacct/wiki/</u>
 - Wiki: docs, implementation notes, ecosystem, etc.



Thanks! Questions?

Paolo Lucente <paolo@pmacct.net>

http://www.pmacct.net/ | https://github.com/pmacct/pmacct

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Backup slides

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Post-processing RDBMS and reporting (1/2)

Traffic delivered to a BGP peer, per location:

mysql> SELECT peer_as_dst, peer_ip_dst, SUM(bytes), stamp_inserted FROM acct_bgp WHERE peer_as_dst = <peer | customer | IP transit> AND stamp_inserted = < today | last hour | last 5 mins > GROUP BY peer as dst, peer ip dst;

Aggregate AS PATHs to the second hop:

mysql> SELECT SUBSTRING_INDEX(as_path, '.', 2) AS as_path, bytes FROM acct_bgp WHERE local_pref = < IP transit pref> AND stamp_inserted = < today | yesterday | last week > GROUP BY SUBSTRING_INDEX(as_path, '.', 2) ORDER BY SUM(bytes);

Focus peak hour (say, 8pm) data:

mysql> SELECT ... FROM ... WHERE stamp_inserted LIKE '2010-02-% 20:00:00'

Post-processing RDBMS and reporting (2/2)

Traffic breakdown, ie. top N grouping BGP peers of the same kind (ie. peers, customers, transit):

mysql> SELECT ... FROM ... WHERE ...

...

```
local_pref = <<pre>ref | customer | IP transit> pref>
```

 Download traffic matrix (or a subset of it) to 3rd party backbone planning/traffic engineering application (ie. Cariden, Wandl, etc.):

```
mysql> SELECT peer_ip_src, peer_ip_dst, bytes, stamp_inserted
    FROM acct_bgp
    WHERE [ peer_ip_src = <location A> AND
        peer_ip_dst = <location Z> AND ... ]
        stamp_inserted = < today | last hour | last 5 mins >
        GROUP BY peer_ip_src, peer_ip_dst;
```